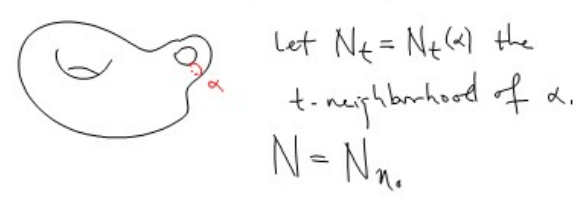


Lecture 15: Proving Separation

$$\forall S \exists \eta_0. \forall \alpha \in \mathcal{Q}(S) \exists \beta \in \mathcal{C}(S) \\ d_g(\alpha, \beta) \geq \eta_0.$$

Recall: Any spine has length $\geq \eta_1$
 Let $\eta_0 = \frac{\eta_1}{100 + 2n}$, $n = \text{genus}(S) + |\mathbb{P}|$.

Also, picked $\alpha \in \mathcal{S}(S)$ with length arb. close to systole length.



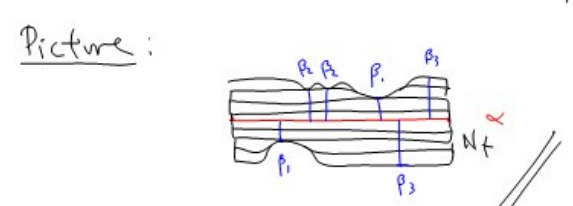
For a contradiction suppose

Claim: \exists arcs $\beta_1, \beta_2, \dots, \beta_n$ s.t.
 * $\beta_i \subseteq N$, $\text{Length}(\beta_i) \leq 2\eta_0$, $\partial\beta_i \subseteq \alpha$
 $\sigma = \alpha \cup \beta_1 \cup \beta_2 \cup \dots \cup \beta_n$ is a spine.

If not: If there is a nontrivial region of $S \setminus N$ then we are done.

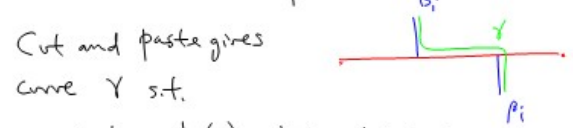
[If we have found separated curves α and some $\beta \subseteq S \setminus N$]

For t small N_t is an annulus. As $t \rightarrow \eta_0$ N_t gathers the topology of S . Every time N_t "bumps" find new arc β_i .



$$\eta_1 \leq \text{Length}(\sigma) \leq \text{Length}(\alpha) + 2n\eta_0 \\ \therefore 100\eta_0 \leq \text{Length}(\alpha).$$

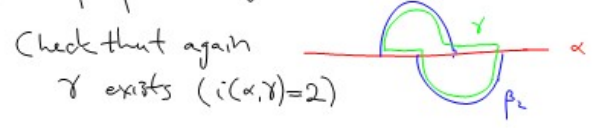
Cases: ① Some β_i meets both sides of α



$$\text{Length}(\gamma) \leq \frac{1}{2}(\text{Length}(\alpha) + 2\eta_0) \\ \leq \text{Length}(\alpha) - 48\eta_0$$

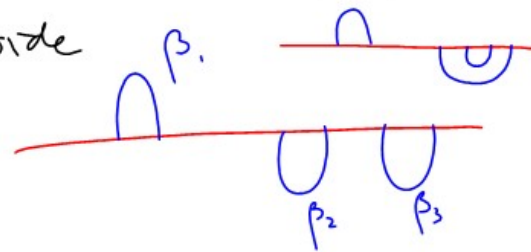
* $i(\alpha, \gamma) = 1$. So γ is ess, non peripheral
 * $\Rightarrow \text{Length}(\gamma) < \text{systole length}$ *

② $\exists \beta_1, \beta_2$ linking waves.



[Check details]

③ $\exists \beta_1$ on one side unlinked from β_2, β_3
on the other side

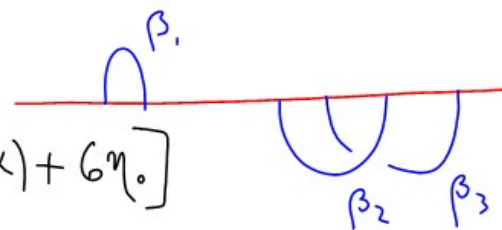


Again find γ

[This time

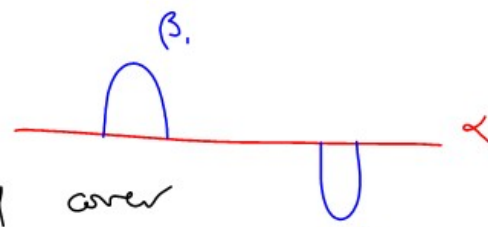
or

$$\text{length}(\gamma) \leq \frac{2}{3} \text{length}(\alpha) + 6\eta_0]$$

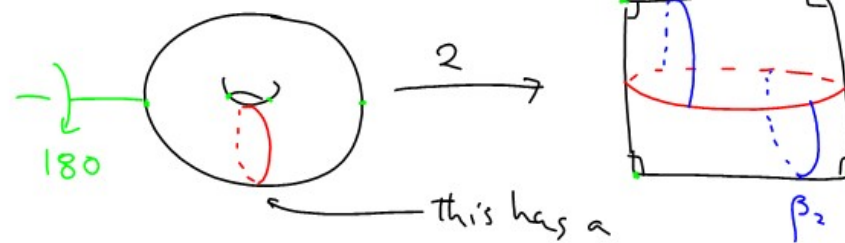


④ $m=2$, no linking, α separating

Have $S \cong S_{0,4}$.



So take double branched cover
and reduce to analysis of $S_1 \cong \mathbb{T}^2$
dealt with in case ① above.



this has a wide annulus, and width goes down by at most a factor of 2 under hyperelliptic.

Exercise: Check cases ②, ③ carefully.

Separation



Annulus inequality: If A is a Riemannian annulus with width w and length l

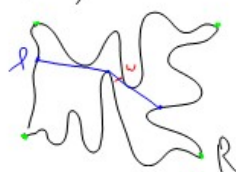


then $\text{Area}(A) \geq w \cdot l$.

[l = length of systole
= inf length of core curves]

Pf: Rectangle inequality

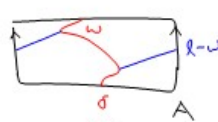
$\text{Area}(R) \geq w \cdot l$



Let A be any annulus, let $\delta \subseteq A, \partial \delta \in \mathbb{Z}A$
s.t. length of δ is as close as needed to w .

Cut A along δ to find

$\text{Area}(A) \geq w \cdot (l - w)$



[via Rectangle inequality for $A \setminus \delta$.]

Let A_n be the n -fold cover of A

Easy: $\text{Area}(A_n) = n \cdot \text{Area}(A)$

Exercise: $w(A_n) = w(A)$ [Easy]

Exercise: $l(A_n) = n \cdot l(A)$

$$n \cdot \text{Area}(A) = \text{Area}(A_n) \geq w(A_n) \cdot (l(A_n) - w(A_n))$$

$$\geq w \cdot (nl - w)$$

$\therefore \text{Area}(A) \geq w \cdot l - \frac{w^2}{n} \xrightarrow{n \rightarrow \infty} w \cdot l$

Corollary: $\forall \gamma \in \mathcal{C}(S), \forall g \in \mathcal{Q}'(S)$

$w_g(\gamma) \cdot l_g(\gamma) \leq 1.$

Pf: A is an (almost) optimal annulus for γ

so $w_g(\gamma) = w_g(A)$, $l_g(\gamma) \leq l_g(A)$
optimal $\forall A$

So: $w_g(\gamma) \cdot l_g(\gamma) \leq w_g(A) \cdot l_g(A) \leq \text{Area}(A)$

$\leq \text{Area}(S) = 1. //$

Wide Curves Define: $W(g, \epsilon) =$

$\{ \gamma \in \mathcal{C}(S) \mid w_g(\gamma) \geq \epsilon \}$

Wide annuli \Rightarrow if $\epsilon \leq \epsilon_0$ then $W(g, \epsilon) \neq \emptyset$.

Set $R_0 = \frac{1}{\epsilon_0}, R = \frac{1}{\epsilon}$.

Thus: By corollary $W(g, \epsilon) \subseteq L(g, R)$
of Annulus inequality.

Observe: If $\delta, \gamma \in \mathcal{C}(S), g \in \mathcal{Q}'(S)$

$w_g(\delta) \cdot i(\delta, \gamma) \leq l_g(\gamma)$. [Exercise]