

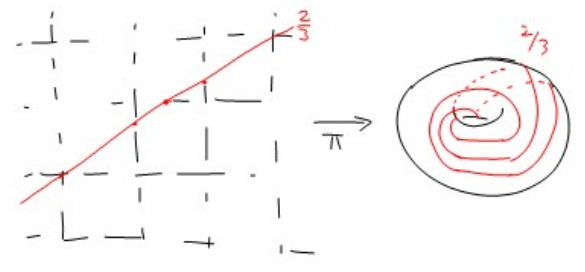
Ⓟ Consider  $\mathbb{T}^2 = S_1$ . Note

$\mathbb{T}^2 = \mathbb{R}^2 / \mathbb{Z}^2$ . Let  $\pi: \mathbb{R}^2 \rightarrow \mathbb{T}^2$

Exercise: ① If  $L \subset \mathbb{R}^2$  is a line of rational slope then  $\pi(L) \in \mathcal{C}^*(\mathbb{T}^2)$

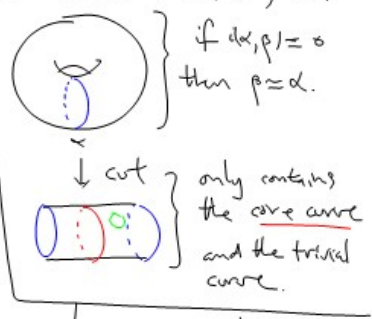
② if  $L, L'$  have same rational slope then  $\pi(L) \cong \pi(L')$

③ [Harder] Every  $\alpha \in \mathcal{C}^*(\mathbb{T}^2)$  is isotopic to  $\pi(L)$ ,  $L$  rational. Say the slope of  $\alpha$  is the slope of  $L$ .



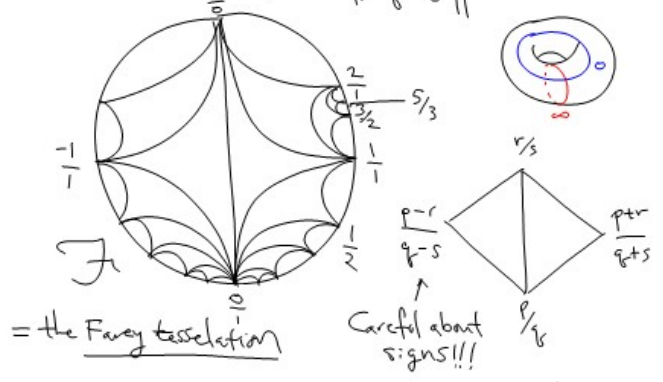
NB: With previous definition  $\mathcal{C}(\mathbb{T}^2)$  has no edges.

Fix this:  
 $\Delta \subseteq \mathcal{C}^*(\mathbb{T}^2)$   
 is a simplex iff



$\forall \alpha, \beta \in \Delta, i(\alpha, \beta) \equiv 1$ . Def:  $i(\alpha, \beta) \equiv 1$  iff  $\alpha, \beta$  are Farey neighbors

Exercise: if  $\alpha, \beta$  have slopes  $\frac{p}{r}, \frac{r}{s}$   
 Then  $i(\alpha, \beta) = \begin{vmatrix} p & r \\ r & s \end{vmatrix}$



Rule: Farey addition only "works" for Farey neighbors. Farey addition

Exercise: Prove that this picture is correct.

- ①  $|\Delta| \leq 3$ . (at most triangles.)
- ② Every vertex meets as many edges
- ③ " edge " exactly 2 triangles.
- ④ Every edge separates  $\mathcal{F}$ .
- ⑤  $\mathcal{F}$  is connected.

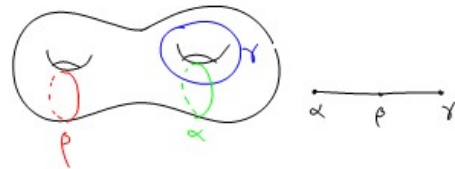
Exercise:  $\mathcal{C}(S_1) \cong \mathcal{C}(S_{1,1}) \cong \mathcal{C}(S_{0,4}) = \mathcal{F}$   
 $i(\alpha, \beta) = 1$        $i(\alpha, \beta) = 2$ .

Exercise: Draw the slope  $1/2$ . [Draw  $3/2$  if too easy]

Ⓓ Metric on  $\mathcal{C}(S)$ . Suppose  $\alpha, \beta \in \mathcal{C}(S)$

Define  $d_S(\alpha, \beta) =$

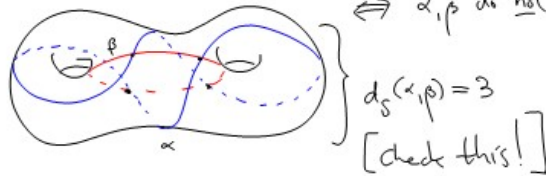
$$\min \left\{ |P| \mid P \text{ an edge path from } \alpha \text{ to } \beta \right\}$$



Review:  $d_S(\alpha, \beta) \leq 1 \iff \alpha, \beta$  disjoint  
( $i(\alpha, \beta) = 0$ )

$d_S(\alpha, \beta) \leq 2 \iff \exists \delta$  disjoint from both.

$\iff \alpha, \beta$  do not fill  $S$ .



Lemma [Hempel] [Barditch, MM, Lickorish, Ivanov, Dehn... various versions]

If  $i(\alpha, \beta) \neq 0$  then

$$d_S(\alpha, \beta) \leq 2 \cdot \log_2(i(\alpha, \beta)) + 2 //$$

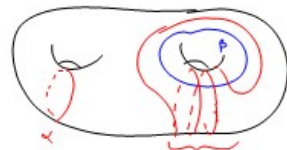
Remark: Sharper version for  $\mathcal{F} = \mathcal{C}(\mathbb{T}^2)$

if  $\alpha, \beta \in \mathcal{F}^\circ$  and  $\alpha \neq \beta$

$$d_{\mathcal{F}}(\alpha, \beta) \leq \log_2(i(\alpha, \beta)) + 1$$

[Pf is easier than pf of Hempel's Lemma]

Neither of these inequalities can be "reversed".



Note:  $d_S(\beta, \beta_n) = 2$  (if  $n \neq 0$ )

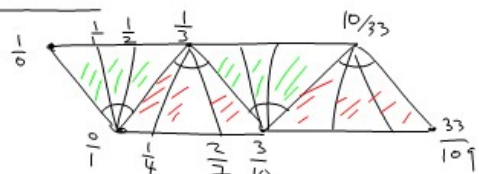
but  $i(\beta, \beta_n) = n$ .

Good Question: Can Hempel's inequality be improved when  $i(\alpha, \beta)$  is large?

Remark: Due to existence of pseudo-Anosov maps, the log must appear.

Corollary:  $\mathcal{C}(S)$  is connected.

Example in  $\mathcal{F}$



[Note  $S_{n+1} = 3S_n + S_{n-1}$ ]

Notice:  $\frac{33}{109} = \frac{1}{3 + \frac{1}{3 + \frac{1}{3 + \frac{1}{3}}}}$   $= [0, 3, 3, 3, 3]$

Call:  $\frac{0}{1}, \frac{1}{3}, \frac{3}{10}, \frac{10}{33}$  the pivots Thus  $d_{\mathcal{F}}(\frac{1}{3}, \frac{33}{109}) \leq 5$

① Prove this is an equality ② Show  $d_{\mathcal{F}}(\frac{1}{3}, \frac{33}{109}) =$  length of cont. frac. of  $\frac{33}{109}$ .