

Lecture 7

Def: A **P-sunflower** is a family of P sets with identical pairwise intersection. That is,

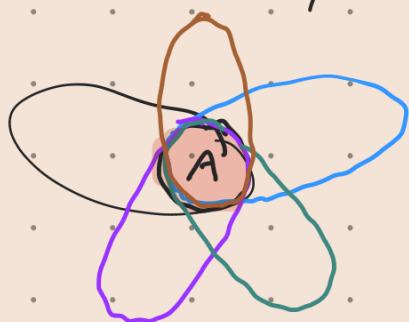
$$\{S_1, \dots, S_P\} \text{ s.t. } \forall i, j \in \binom{[P]}{2} \Rightarrow S_i \cap S_j = A.$$

We call A the **core** of the sunflower

Note that A could be empty.

Here $P = \#$ of petals.

• $\text{core} \neq \emptyset$



• $\text{core} = \emptyset$



• k -graphs = k -uniform hypergraphs = k -sets system

Conj (Erdős-Rado sunflower conj)

Let $P \geq 3$. Then $\exists C = C(P)$ s.t.

\forall k -graph \mathcal{F} without a P -sunflower has

$$\leq C^k \text{ edges.}$$

• Alweiss-Lovett-Wu-Zhang: $(\log k \cdot (P \log \log k)^{O(1)})^k$

• Rao, using Shannon noiseless coding thm,
 $(O(1) \cdot P \log(Pk))^k$

• Tao, entropy pf.

Notations • X set, write $X_\delta = \text{Binom}(X, \delta)$

is binomial random subsets of X w/ each element included w/ prob. δ , indep of others.

• A set, $W \sim A$

W uniform sampled from A .

• W, Z r.v. write

$W \sim Z$ the same distribution.

§ 2.9.1 Reduction to spread hypergraphs
via induction.

Take $R \geq 1$.

Def: (**Spread**) A k -graph \mathcal{H} is R -spread if

$$\forall Z \neq \emptyset, \quad d_{\mathcal{H}}(Z) = |\{S \in \mathcal{H} : S \supseteq Z\}| \leq \frac{|\mathcal{H}|}{R|Z|}$$

• Reduction: We may consider spread hypergraphs

Suppose we want to show:

sunflower-free $\Rightarrow \leq R^k$ edges.

If \mathcal{H} is not R -spread.

$$\Rightarrow \exists Z \text{ s.t. } d(Z) \geq \frac{e(\mathcal{H})}{R|Z|}$$

Now if $e(\mathcal{H}) \geq R^k \Rightarrow d(z) \geq R^{k-|z|}$

consider link of z : $(k-|z|)$ -uniform hypergraph
with edge set $\{s_i \setminus z : s_i \in \mathcal{H}, s_i \supseteq z\}$

Induct. on uniformity \checkmark . 

• We will now find in a spread hypergraph

P pairwise disjoint sets.



Def: A random set A is R -spread, $R > 1$,

if $\Pr(z \subseteq A) \leq R^{-|z|} \quad \forall z$.

A hypergraph \mathcal{H} is R -spread if the uniform random edge from \mathcal{H} is R -spread.

§ 2.9.2. Incorporate spreadness in entropy

• Maximality $\Rightarrow X$ takes values in a set S_Y .

that depends on Y , then $H(X|Y) \leq \mathbb{E}_Y \log |S_Y|$.

\Rightarrow (6) ... $H(B|A) \leq \mathbb{E}|A|$, \forall random $B \subseteq$ random A

- Mutual info. of two r.v. X & Y .

$$\begin{aligned} I(X;Y) &= H(X) + H(Y) - H(X,Y) \\ &= H(X) - H(X|Y) = H(Y) - H(Y|X) \\ &= H(X,Y) - H(X|Y) - H(Y|X) \end{aligned}$$

Monotonicity

- if $Y = f(X) \Rightarrow I(X;Y) = H(Y)$

indep.

- if X, Y indep. $\Rightarrow I(X;Y) = 0$

- Let X be a r.v., an empirical seq X_1, X_2, \dots for X that take values from the same set as X and converging to X in distribution, that is,

$$n \sim [N], \quad \Pr(X_n = a) = \Pr(X = a) + o_N(1)$$

- Mutual info between a random set & its random subset.

Consider $A \sim \mathcal{Z}^{[n]}$ & a random subset $B \subseteq A$.

$$\Rightarrow I(A;B) = H(A) - H(A|B)$$

$$= \mathbb{E}_b (H(A) - H(A|B=b))$$

$$\stackrel{\text{max.}}{=} \mathbb{E}_b (n - (n-b)) = \mathbb{E}_b b = \mathbb{E}|B|$$

- Spread set shares a lot more info. w/ its random subset.

Lem 2.20 Let A be a random R -spread set w/ $|R| > 1$.

If A is uniformly chosen, then \forall random $B \subseteq A$

$$\Rightarrow I(A; B) \geq \log R \cdot \mathbb{E}|B|$$

Rmk: We will use this lem in contrapositive way.

WTS $\mathbb{E}|B|$ small \Rightarrow enough to upp. bd. $I(A; B)$

Lem 2.21 Let $W \sim \binom{X}{\delta|X|}$, $0 < \delta < 1$

- \forall random $B \subseteq W$, $H(W) - H(W|B) \geq \log \frac{1}{\delta} \cdot \mathbb{E}|B|$
- (Absorption) \forall random $B \subseteq X$,

$$H(W \cup B) \leq H(W) + 1 + (1 + \log \frac{1}{\delta}) \cdot \mathbb{E}|B| \dots (7)$$

- Relative product. (X, X') in which X' is a cond. indep. copy of X subject to certain constraint $f(X) = f(X')$ for a given f .

$$\Rightarrow H(X) = H(X, f(X)) = H(X|f(X)) + H(f(X))$$

only relevant info to X from X' is $f(X) = f(X')$ \Downarrow

$$= H(X|X') + H(f(X)) \dots (8)$$

Lem 2.22 Let $R > 1$, $0 < \delta < 1$, A be a random R -spread subset of X , and $V \sim X_\delta$ indep. of A .

Then \exists another random $A' \sim A$ of X s.t.

$$A' \setminus V \subseteq A \quad \text{and} \quad \mathbb{E}|A' \setminus V| \leq \frac{4 + \log \frac{1}{\delta}}{\log R} \cdot \mathbb{E}|A|$$

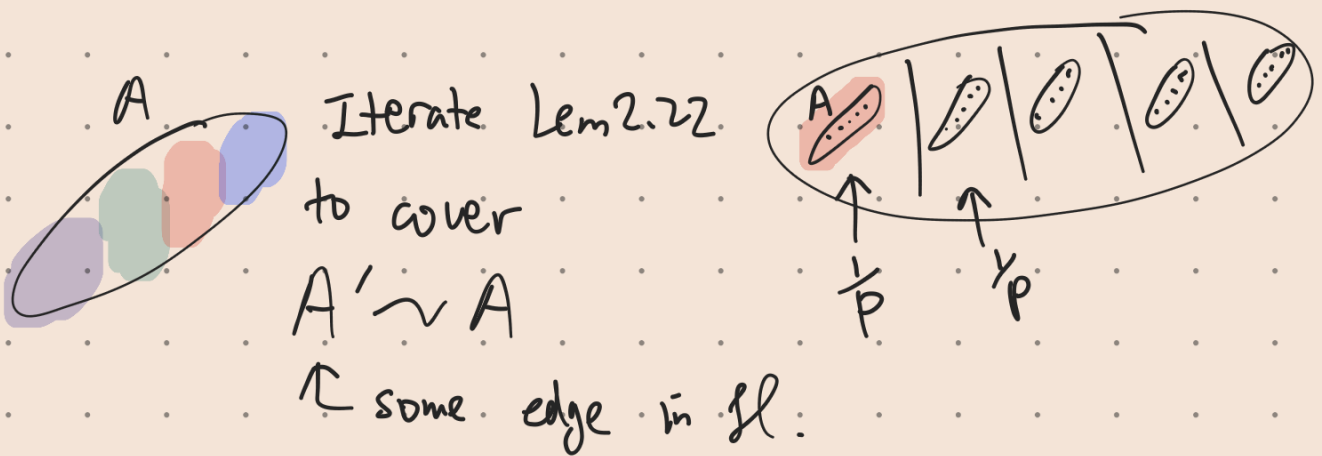
Proof: $\mathbb{E}|A' \setminus V| = (1 - \delta) \mathbb{E}|A|$, by passing to a $A' \sim A$ we can improve $(1 - o(1))$ to $o(1)$ factor (if $R \gg \frac{1}{\delta}$)

(Sketch)

Pf Lem 2.22 \Rightarrow upp bd. $\leq R^k$, where $R = (P \log(Pk))^{O(k)}$

• \mathcal{H} , R -spread

Sample A unif from $\mathcal{H} \Rightarrow A$ R -spread.



I.e. iterate Lem 2.22 so that $\mathbb{E}|A' \setminus V| < \frac{1}{p}$

Markov $\Rightarrow \Pr(|A' \setminus V| \geq 1) \leq \mathbb{E}|A' \setminus V| < \frac{1}{p}$

then union bd. \square

Pf (Lem 2.22). If A is empty, take $A' = A$. \checkmark .

May assume $\mathbb{E}|A| \geq 1$..

• View $X = [|X|]$, take $N_1, N_2 \gg |X|$

• Pass binom. model \Rightarrow unif. model.

$$V \sim X_{\delta} \Rightarrow W \sim \binom{[N_2]}{\delta N_2}$$

Take empirical seq of A : A_1, A_2, \dots

$n \sim [N_1] \Rightarrow$ random set A_n

$$\mathbb{E}|A_n| \geq \mathbb{E}|A| - o(1) \geq 1 - o(1)$$

Note that $W \cap [|X|] \xrightarrow{\text{dist.}} V$

This, as $N_1, N_2 \rightarrow \infty$, suffices to show

$$\exists n' \sim n \quad \mathbb{E}|A_{n'} \setminus W| \leq \frac{4 + \log \frac{1}{\delta}}{\log R} \mathbb{E}|A_n| + o(1)$$

Key idea Take a cond. indep copy (n', W') of

(n, W) Subject to

$$(9) \dots A_n \cup W' = A_n \cup W \Rightarrow A_{n'} \setminus W \subseteq A_n \cap A_{n'}$$

Thus, enough to show

$$\mathbb{E}|A_n \cap A_{n'}| \leq \frac{4 + \log \frac{1}{\delta}}{\log R} \mathbb{E}|A_n| + o(1)$$

By Lem 2.20 w./ $(A, B) = (A_n, A_n \cap A_{n'})$

Suffices: $H(n) - H(n | A_n \cap A_{n'}) \leq (4 + \log \frac{1}{\delta}) \mathbb{E}|A_n| + o(1)$
..... (10)

• bd $H(n)$ using $H(n, W)$

use (8) w./ $(n, W), (n', W'), A_n \cup W = (X, X', f(X))$

$$H(n, W) = H(n, W | n', W') + H(A_n \cup W)$$

$H(n) + H(W)$ ~~ind. //~~ (7) $\leq H(n, W | n', W') + H(W) + (2 + \log \frac{1}{\delta}) \mathbb{E}|A_n| + o(1)$

$\Rightarrow H(n) \leq H(n, W | n', W') + \underbrace{\dots}_{(11)}$

• Relate $H(n, W | n', W')$ to $A_n \cap A_{n'}$

$$\begin{aligned} H(n, W | n', W') &\leq \underbrace{H(A_n \cap A_{n'} | n', W')}_{\text{choose } A_n \cap A_{n'}} \\ &\quad + H(n | A_n \cap A_{n'}, n', W') \quad \text{choose } A_n \\ &\quad + \underbrace{H(W | n, A_n \cap A_{n'}, n', W')}_{\text{choose } W} \leq \mathbb{E}|A_n| \end{aligned}$$

1st & 3rd term can be bdd by (6) \uparrow

$\Rightarrow H(n, W | n', W') \leq 2 \mathbb{E}|A_n| + H(n | A_n \cap A_{n'})$

plug this into (11) \square